

# The higher sharp

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Equivalently,

- ▶  $A \subseteq \mathbb{R}$  is  $\Sigma_1^1$  iff there is a Borel set  $B \subseteq \mathbb{R}^2$  such that  $x \in A \leftrightarrow \exists y (x, y) \in B$ .

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$A$  is **projective** iff  $A$  is  $\Sigma_n^1$  for some  $n < \omega$ .



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- ▶ (Woodin) Suppose there are  $\omega$  Woodin cardinals below a measurable cardinal, then  $L(\mathbb{R}) \models \text{AD}$ .

In this talk, we assume  $\text{ZFC} + L(\mathbb{R}) \models \text{AD}$ .

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## Question

*What about the effective version?*

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DST way of thinking: Suppose  $\emptyset \neq A = p[T]$ .

- ▶  $(x, \vec{\alpha}) \in [T]$  is the **leftmost branch** of  $T$  iff whenever  $(y, \vec{\beta}) \in [T]$ ,  $(x(0), \alpha_0, x(1), \alpha_1, \dots)$  is lexicographically smaller than  $(y(0), \beta_0, y(1), \beta_1, \dots)$ .



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- ▶ If  $(x, \vec{\alpha}) \in [T]$  is the leftmost branch of  $T$ , then  $x$  is the **leftmost real** associated to  $T$ .

# Effective codes of ordinals (1)

$\Sigma_2^1$  sets are “effectively”  $\omega_1$ -Suslin.

Codes of ordinals up to  $\omega_1$ :

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Is there a canonical coding system for ordinals smaller than  $\Theta^{UB}$ ?

Next level:  $\Sigma_3^1$  sets are “effectively”  $\omega_\omega$ -Suslin.

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- ▶  $L$  is Skolem generated by  $\{c_\alpha : \alpha < \infty\}$ ,
- ▶ for every first-order formula  $\varphi$ , every  $\alpha_1 < \dots < \alpha_n$ ,  
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$$L \models \varphi(c_{\alpha_1}, \dots, c_{\alpha_n}) \leftrightarrow \varphi(c_{\beta_1}, \dots, c_{\beta_n}).$$

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$(c_\alpha : \alpha < \infty)$  is called the class of **Silver indiscernibles** for  $L$ . In particular, every uncountable cardinal is a Silver indiscernible for  $L$ .

$$0^\# = \bigoplus_{n < \omega} \{ \ulcorner \varphi \urcorner : \varphi \text{ has } n \text{ free variables, } L \models \varphi(\aleph_1, \dots, \aleph_n) \}.$$

$0^\#$  is the unique wellfounded remarkable EM blueprint.



# The EM blueprint formulation of sharps

## Definition

An **EM blueprint** is a complete consistent theory in the language  $\{\dot{\epsilon}, \dot{c}_n : n < \omega\}$  that extends the following axioms:

- ▶  $ZFC + V = L$ . In particular, there is a definable wellordering of the universe.
- ▶  $c_1 < c_2$
- ▶ For every formula  $\varphi$ , every increasing tuple  $k_1 < \dots < k_m$ ,  $\varphi(\dot{c}_1, \dots, \dot{c}_m) \rightarrow \varphi(\dot{c}_{k_1}, \dots, \dot{c}_{k_m})$  is an axiom

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An EM blueprint  $\Gamma$  is **remarkable** iff  $\Gamma$  contains the remarkability axioms:

$$\tau(\dot{c}_1, \dots, \dot{c}_{m+n}) < \dot{c}_{m+1} \rightarrow \tau(\dot{c}_1, \dots, \dot{c}_{m+n}) = \tau(\dot{c}_1, \dots, \dot{c}_m, \dot{c}_{m+n+1}, \dots, \dot{c}_{m+n+1})$$

holds for every Skolem term  $\tau$ .

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Suppose  $\Gamma$  is an EM blueprint. For every linear ordering  $<^*$  on a set  $W$ , there is (up to isomorphism) a unique model  $\mathcal{M}(\Gamma, <^*)$  and points  $\{b_v : v \in W\} \subseteq \mathcal{M}(\Gamma, <^*)$  such that

$$\mathcal{M}(\Gamma, <^*) \models \varphi(b_{v_1}, \dots, b_{v_n})$$

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**Wellfoundedness is  $\Pi_2^1$ .**

- ▶  $0^\#$  is the unique wellfounded remarkable EM blueprint
- ▶  $x^\#$  is the unique wellfounded remarkable EM blueprint over  $x$ .

## Effective codes of ordinals (2)

Codes of ordinals up to  $u_\omega$ :

$\gamma$  is a **uniform indiscernible** iff  $\gamma$  is a Silver indiscernible for  $L[x]$  for every  $x \in \mathbb{R}$ . The uniform indiscernibles are enumerated  $(u_\alpha : 1 \leq \alpha < \infty)$ .

$$u_1 = \omega_1. \quad u_\omega = \sup_{n < \omega} u_n.$$

A **sharp code** is pair  $(\ulcorner \tau \urcorner, x^\#)$  where  $x \in \mathbb{R}$ ,  $\tau$  is an  $n$ -ary Skolem term for some  $n$ , and  $\tau^{L[x]}(u_1, \dots, u_n)$  is an ordinal.



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- ▶ Every ordinal in  $u_\omega$  is represented by a sharp code.

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$\gamma$  is a **uniform indiscernible** iff  $\gamma$  is a Silver indiscernible for  $L[x]$  for every  $x \in \mathbb{R}$ . The uniform indiscernibles are enumerated  $(u_\alpha : 1 \leq \alpha < \infty)$ .

$$u_1 = \omega_1. \quad u_\omega = \sup_{n < \omega} u_n.$$

A **sharp code** is pair  $(\ulcorner \tau \urcorner, x^\#)$  where  $x \in \mathbb{R}$ ,  $\tau$  is an  $n$ -ary Skolem term for some  $n$ , and  $\tau^{L[x]}(u_1, \dots, u_n)$  is an ordinal. If  $(\ulcorner \tau \urcorner, x^\#)$  is an  $n$ -ary sharp code, then it codes an ordinal in  $u_{n+1}$ :

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- ▶ The set of sharp codes is  $\Pi_2^1$ .
- ▶ The relation “ $v, w$  are sharp codes  $\wedge |v| = |w|$ ” is  $\Delta_3^1$ .

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- ▶  $u_1 = \omega_1, u_2, \dots, u_n, \dots$  list the first  $\omega$  uncountable projective wellordered cardinals.
- ▶ The “projective successor” of  $u_\omega$  is  $\delta_3^1$ .
- ▶ IMT provides a deep insight into the structural theory related to the projective well-ordered cardinals.



# The Martin-Solovay tree projecting to $\Pi_2^1$

## Theorem (Martin-Solovay)

*If  $A$  is  $\Pi_2^1$ , then  $A = p[T]$  for some  $T$  on  $\omega \times u_\omega$  such that  $T$  is  $\Delta_3^1$  in the sharp codes.*

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## Definition of the Martin-Solovay tree:

If  $\sigma : \{1, \dots, m\} \rightarrow \{1, \dots, n\}$  is an order preserving map, then  $\sigma$  induces  $j^\sigma : u_{m+1} \rightarrow u_{n+1}$  by

$$j^\sigma(\tau^{L[x]}(u_1, \dots, u_m)) = \tau^{L[x]}(u_{\sigma(1)}, \dots, u_{\sigma(m)}).$$

# The Martin-Solovay tree projecting to $\aleph_2^1$

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$$\tau_m(\dot{c}_{\sigma(1)}, \dots, \dot{c}_{\sigma(k_m)}) = \tau_n(\dot{c}_1, \dots, \dot{c}_{k_n})$$

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- ▶ (Martin-Solovay)  $p[T_2] = \{x^\# : x \in \mathbb{R}\}$ .
- ▶ It is easy to define  $\widehat{T}_2$ , a variant of  $T_2$ , such that  $p[\widehat{T}_2]$  is a good universal  $\Pi_2^1$  set.

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$\widehat{T}_2$  is a tree on  $\omega \times u_\omega$ . It is  $\Delta_3^1$  in the codes.  $p[\widehat{T}_2]$  is a good universal  $\Pi_2^1$  set.



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## Theorem

Assume  $\Delta_2^1$ -determinacy. Then  $y_3^0 \equiv_m M_1^\#$  (many-one equivalence).

## Effective codes of ordinals (3)

Codes of ordinals up to  $\delta_3^1$ :

Recall: Under AD,  $u_n = \aleph_n$ ,  $\lambda_3 = u_\omega = \aleph_\omega$ ,  $\delta_3^1 = \aleph_{\omega+1}$ .

- ▶  $x \in \text{WO}_{(3)}$  iff  $x$  codes a wellordering of  $u_\omega$  via Kunen's coding of subsets of  $u_\omega$ . (under AD, there is a surjection  $\pi : \mathbb{R} \rightarrow \mathcal{P}(u_\omega)$  such that the relation  $\alpha \in \pi(x)$  is  $\Delta_3^1$ )

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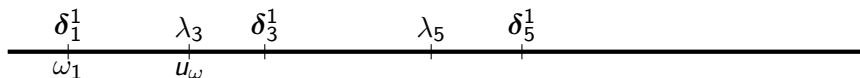
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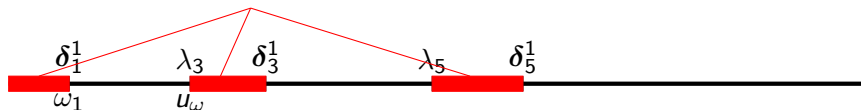
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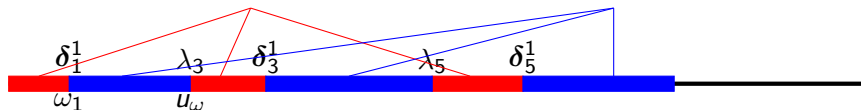
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The model theoretic form of  $\Pi_1^1$  subsets of  $\mathbb{R}$

The following are equivalent:

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**Higher Kleene's O:**  $(\ulcorner \varphi \urcorner, \alpha) \in \mathcal{O}^{T_2}$  iff  $\varphi$  is a  $\Sigma_1$  formula with one free variable,  $\alpha < u_\omega$  and  $L_{\kappa_3}[T_2] \models \varphi(\alpha)$ .

## The level-2 sharp

$\mathcal{O}^{T_2} \notin L_{\kappa_3}[T_2]$ . For each  $n < \omega$ ,  $\mathcal{O}^{T_2} \cap (\omega \times u_n) \in L_{\kappa_3}[T_2]$ .

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# The level-2 sharp

$\mathcal{O}^{T_2} \notin L_{\kappa_3}[T_2]$ . For each  $n < \omega$ ,  $\mathcal{O}^{T_2} \cap (\omega \times u_n) \in L_{\kappa_3}[T_2]$ .

## Definition

- ▶  $(0^{2\#})_n = \{(\varphi, \psi) : \varphi, \psi \text{ are } \Sigma_1 \text{ formulas, } \exists \alpha < u_n$   
 $((\varphi, \alpha) \notin \mathcal{O}^{T_2} \wedge \forall \eta < \alpha (\psi, \eta) \in \mathcal{O}^{T_2})\}$ .
- ▶  $\mathcal{O}^{T_2} \cap (\omega \times u_n)$  is squeezed into the real  $(0^{2\#})_n$  by applying a Boolean operation on the second coordinate.
- ▶  $0^{2\#} = \bigoplus_{n < \omega} (0^{2\#})_n$ .
- ▶ For  $x \in \mathbb{R}$ ,  $x^{2\#}$  is the relativized version.

By the previous slide,  $0^{2\#} \equiv_m y_3^0$ . It is the model theoretic definition of  $y_3^0$ .

## Theorem

$0^{2\#} \equiv_m M_1^\#$ .  $x^{2\#} \equiv_m M_1^\#(x)$ , the many-one reductions independent of  $x$ .

## $\Pi_3^1$ sets, the structure of $L[T_3]$

### Definition

$T_3$  is the tree on  $\omega \times \delta_3^1$  associated to a  $\Pi_3^1$ -scale on a good universal  $\Pi_3^1$  set.

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1. (Becker-Kechris)  $L[T_3]$  is invariant of the choice of  $T_3$ .
2. (Steel)  $L[T_3]$  is a mouse. Let  $j : M_2^\# \rightarrow M_{2,\infty}^\#$  be the direct limit map of all countable iterates of  $M_2^\#$ ,  $\delta_\infty$  be the least Woodin of  $M_{2,\infty}^\#$ , then
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Intuition:  $(L[T_3], \text{level-3 indiscernibles}) \sim M_2^\#$ .

# Level-3 EM blueprint

## The role of level-3 EM blueprints

The (level-1) EM blueprint definition of sharps leads to

- ▶ Effective wellordered cardinals  $(u_n : n \leq \omega)$ ,  $u_n = \aleph_n^{L(\mathbb{R})}$ ,
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## Level-3 indiscernibles

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- ▶  $\mathcal{L}^R$  is the language  $\{\dot{c}, \dot{c}_s : s \in \text{dom}(R)\}$ .
- ▶  $\Gamma(R)$  is a complete consistent  $\mathcal{L}^R$  theory containing the following axioms:
  - ▶  $ZFC + V = K +$  “I am a 2-small mouse”.
  - ▶ ( $\varphi_\Sigma, \varphi_\Pi$  are fixed formulas)  $\varphi_\Sigma, \varphi_\Pi$  define class trees  $T_\Sigma, T_\Pi$  resp. In any set generic extension,  $\check{T}_\Sigma, \check{T}_\Pi$  project to the good universal  $\Sigma_3^1$  set and its complement.
  - ▶  $V$  is closed under the  $M_1^\#$  operator, as certified by  $T_\Pi$ .
  - ▶ ... (The order relation of those  $c_s$ 's) ...

$\mathcal{M}(\Gamma, R)$  is (up to isomorphism) the unique model of  $\Gamma(R)$  that is Skolem generated by  $(\dot{c}_s^{\mathcal{M}(\Gamma, R)} : s \in \mathbb{R})$ .  $\mathcal{M}(\Gamma, R)$  is called the **level-3 EM model** associated to  $\Gamma$  and  $R$ .



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- ▶  $\Gamma$  is **coherent**. In particular, if  $R$  is a subfunction of  $R'$ , then  $\Gamma(R) \subseteq \Gamma(R')$ , so that  $\mathcal{M}(\Gamma, R)$  elementarily embeds into  $\mathcal{M}(\Gamma, R')$  in a unique way.

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- ▶ **Remarkability**. Generalization of the level-1 remarkability axiom.
- ▶ **Iterability**. If a tower of level-3 trees  $(R_n : n < \omega)$  is  $\Pi_3^1$ -wellfounded, then the direct limit of  $(\mathcal{M}(\Gamma, R_n) : n < \omega)$  is a  $\Pi_3^1$ -iterable mouse.

**Iterability is  $\Pi_4^1$ .**

**Theorem**

*Assume  $\Pi_3^1$ -determinacy. Then there is a unique iterable remarkable level-3 EM blueprint.*

# The level-3 sharp

## Definition

Assume  $\Pi_3^1$ -determinacy.

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The proof uses the equivalence  $x^{2\#} \equiv_m M_1^\#(x)$ .

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The level-3 EM blueprint definition of level-3 sharps leads to

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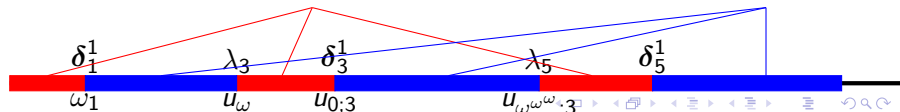
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WO or (higher level) Kunen's coding

(higher level) sharp coding



# The model theoretic form of $\Pi_5^1$ sets

The model theoretic form of  $\Pi_5^1$  subsets of  $u_{\omega^{\omega};3} \times \mathbb{R}$

The following are equivalent:

1.  $A \subseteq u_{\omega^{\omega};3} \times \mathbb{R}$  is  $\Pi_5^1$ .
2. For some  $\Sigma_1$  formula  $\varphi$ ,  $x \in A$  iff  $L_{\kappa_5^x}[T_4, x] \models \varphi(T_4, \alpha, x)$ .

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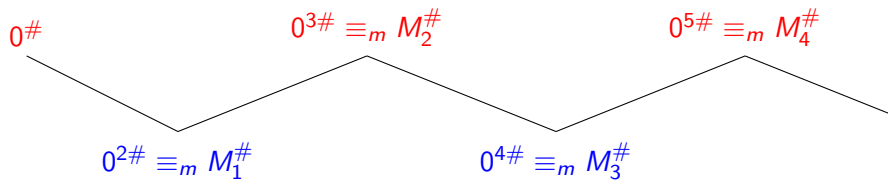
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$\mathcal{O}^{4\#}$  is obtained by squeezing  $\mathcal{O}^{T_4}$  into a real. It is many-one equivalent to  $M_3^\#$ .

# Continue!

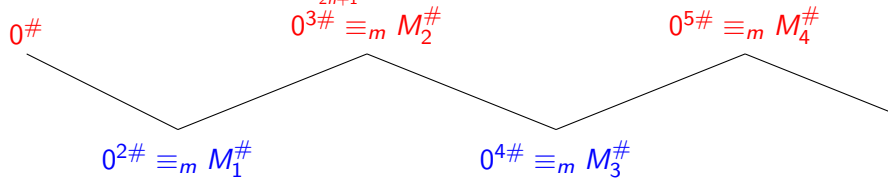


## Question

1. *How to continue this analysis in  $L(\mathbb{R})$  and beyond? I.e., an effective coding of ordinals in  $\Theta^{L(\mathbb{R})}$ , etc.*
2. *Is there a purely inner model theoretic approach? Test question: What is the fine structure of  $L_{\kappa_3}[T_2]$ ?*

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Theory of the mouse  $L_{\delta_{2n+1}^1}[T_{2n+1}]$  with level- $(2n+1)$  indiscernibles

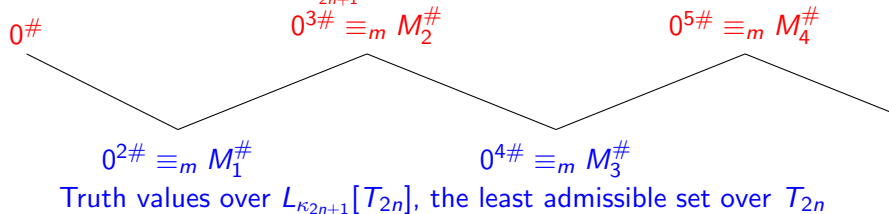


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# Applications

## Theorem

*The following are equivalent.*

1.  $\Pi_{2n+1}^1$  and  $\Pi_{2n+2}^1$  sets are determined.
2. There is a countably iterable  $M_{2n+1}^\#$ .

The  $n = 0$  case is proved by Neeman and Woodin. The equivalence of  $0^{(n+1)\#}$  and  $M_n^\#$  allows the full generalization. (Open for  $n > 0$ ) What if we replace  $(2n + 1, 2n + 2)$  by  $(2n, 2n + 1)$ ?

## The big hope

Every theorem in DST at the level of  $\Sigma_1^1$  under the assumption of sharps will generalize to projective sets under PD. E.g., the classification of thin projective equivalence relations on  $\mathbb{R}$ .

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Thank you!