

Omitting types in infinitary $[0, 1]$ -valued logic

Christopher Eagle
University of Toronto

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Model theory for metric structures

Metric structures: $\langle M, d, (f_i)_{i \in I}, (R_j)_{j \in J}, (c_k)_{k \in K} \rangle$

- ▶ (M, d) metric space of diameter ≤ 1
 - ▶ Not necessarily complete
- ▶ $f_i : M^n \rightarrow M$, uniformly continuous
- ▶ $R_j : M^n \rightarrow [0, 1]$, uniformly continuous
- ▶ $c_k \in M$.

Example

First-order structures, (unit balls of) Banach spaces, operator algebras, ...

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- ▶ Signatures specify moduli of uniform continuity for functions and predicates.
- ▶ Terms defined exactly as in first-order.
- ▶ Atomic formulas:
 - ▶ $d(t_1(\bar{x}), t_2(\bar{x}))$
 - ▶ $R(t_1(\bar{x}), \dots, t_n(\bar{x}))$
 - ▶ Constant formula r for each $r \in \mathbb{Q} \cap [0, 1]$.
- ▶ Formulas:
 - ▶ Atomic formulas
 - ▶ $\phi \rightarrow \psi (= \max\{\phi - \psi, 0\})$
 - ▶ $\exists \epsilon > 0$ Other connectives $\mathbb{Q} \cap [0, 1] \rightarrow [0, 1]$ can be approximated
 - ▶ $\inf_x \phi$ (approximate existential quantifier)

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- ▶ Formula $\phi(x_1, \dots, x_n)$ defines a function on each structure M
 $\phi^M : M^n \rightarrow [0, 1]$.
- ▶ Write $M \models \phi(a_1, \dots, a_n)$ to mean $\phi^M(a_1, \dots, a_n) = 0$.
- ▶ $M \models \phi(\bar{a}) \rightarrow \psi(\bar{b}) \iff \phi^M(\bar{a}) \leq \psi^M(\bar{b})$.
- ▶ This logic is similar to first-order in many ways.
 - ▶ See survey article by Ben Yaacov, Berenstein, Henson, and Usvyatsov.

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Infinitary $[0, 1]$ -valued logic

- ▶ Goal: Extend to an infinitary logic with properties similar to $L_{\omega_1, \omega}$.
 - ▶ Scott sentences for separable structures.
 - ▶ Omitting types for countable fragments.
- ▶ Given formulas ϕ_0, ϕ_1, \dots , allow a formula $\inf_{n < \omega} \phi_n$ (approximate infinitary disjunction).
 - ▶ Call this logic $L_{\omega_1, \omega}[0, 1]$.
 - ▶ If $\phi(x_1, \dots, x_m)$ is an $L_{\omega_1, \omega}[0, 1]$ -formula, then $\phi^M : M^n \rightarrow [0, 1]$ is not necessarily continuous!
 - ▶ Discontinuous formulas can be used to study structures from functional analysis, but care is needed.

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Examples of notions from Banach space theory expressible in $L_{\omega_1, \omega}[0, 1]$ (in signatures expanding that of Banach spaces) include:

- ▶ A sequence (named as constant symbols) is a Schauder basis.
- ▶ Failure of hereditary indecomposability.
- ▶ Failure of reflexivity.
- ▶ Equivalence of norms.
- ▶ ...

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Scott sentences

Theorem

For each separable Banach space X there is an $L_{\omega_1, \omega}[0, 1]$ -sentence σ such that if Y is separable and $Y \models \sigma$, then Y is almost isometric to X .

- ▶ The proof is an analysis of the back-and-forth games introduced by Heinrich and Henson, analogous to Ehrenfeucht-Fraïssé games.
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Omitting Types - Definitions

- ▶ A **fragment** of $L_{\omega_1, \omega}[0, 1]$ is a set of $L_{\omega_1, \omega}[0, 1]$ formulas containing the atomic formulas and closed under substitutions, subformulas, finitary connectives, and (approximate) existential quantification.
- ▶ A **type** of a theory T is a set of formulas which is satisfied in some model of T .
- ▶ A type $\Sigma(\bar{x})$ is **principal** over T if there is a formula $\phi(\bar{x})$ and an $r \in \mathbb{Q} \cap (0, 1)$ such that:
 - ▶ ϕ is consistent with T ,
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Theorem

Let L be a countable fragment of $L_{\omega_1, \omega}[0, 1]$, and T an L -theory. If $\{\Sigma_n(\bar{x}) : n < \omega\}$ is a set of non-principal types of T , then there is a countable model of T omitting each Σ_n .

- ▶ With a slightly modified definition of principal, this implies a version where the resulting structure is separable and complete.
- ▶ Restricted to discrete structures - Keisler
- ▶ Restricted to the finitary fragment - Henson

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Let L be a countable fragment of $L_{\omega_1, \omega}[0, 1]$, and T an L -theory. If $\{\Sigma_n(\bar{x}) : n < \omega\}$ is a set of non-principal types of T , then there is a countable model of T omitting each Σ_n .

- ▶ With a slightly modified definition of principal, this implies a version where the resulting structure is separable and complete.
- ▶ Restricted to discrete structures - Keisler
- ▶ Restricted to the finitary fragment - Henson

Omitting Types - Topology

- ▶ Omitting types is closely related to Baire Category.
- ▶ Goal: Use topological ideas to give a simple proof.
 - ▶ Build on Morley's topological proof from the discrete infinitary setting.
 - ▶ Minimize the role of syntax.

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- ▶ $Str_L = \{M : M \text{ is an } L\text{-structure}\}$
- ▶ Closed classes: $Mod(T) = \{M \in Str_L : M \models T\}$.
- ▶ Str_L is completely regular, not Hausdorff.

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Omitting Types - Topology

- ▶ C - countable set of new constant symbols,
- ▶ $L(C)$ - least countable fragment containing L and C ,
- ▶ \mathcal{W} - subspace of $Str_{L(C)}$ such that $M \models \inf_x \phi$ is witnessed in C .

Proposition

If every closed subspace of \mathcal{W} is Baire, then the Omitting Types Theorem holds.

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Definition

A sequence $\langle \mathcal{U}_n : n < \omega \rangle$ of open covers of a space X is **complete** if the following holds:

- ▶ If: \mathcal{F} is a centred family of closed sets
- ▶ And: for each n there is $F_n \in \mathcal{F}$ and $U_n \in \mathcal{U}_n$ such that $F_n \subseteq U_n$
- ▶ Then: $\bigcap \mathcal{F} \neq \emptyset$.

A completely regular space X is **Čech-complete** if it has a complete sequence of open covers.

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If X is Čech-complete then every closed subspace of X is Baire.

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\mathcal{W} is Čech-complete.

Idea of proof.

- ▶ Enumerate the sentences as $\sigma_0, \sigma_1, \dots$
- ▶ The sets in \mathcal{U}_n prescribe estimates for $\sigma_0, \dots, \sigma_n$.
- ▶ Estimates at stage $n + 1$ refine those from stage n .
- ▶ When estimating $\inf_{k < \omega} \phi_k$, also choose a k and make the same estimate for ϕ_k .
- ▶ Ultrafilter limits of sequences following these open covers exist. A Łoś-type result shows that for these sequences, ultrafilter limits can be computed as ultraproducts.
 - ▶ Similar to the proof of compactness in first-order, but this space is not compact!



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Application - Separable Quotients

Theorem

Let $T : X \rightarrow Y$ be a bounded surjective linear map, where X and Y are Banach spaces and $\text{density}(X) > \text{density}(Y)$. Let L be a “nice” countable fragment of $L_{\omega_1, \omega}[0, 1]$. Then there exist Banach spaces X', Y' , and a bounded surjective linear map $T' : X' \rightarrow Y'$, such that:

- ▶ $\text{density}(X') = \aleph_1$,
- ▶ $\text{density}(Y') = \aleph_0$,
- ▶ $(X, Y, T) \equiv_L (X', Y', T')$.

This extends a result of Ben Yaacov and Iovino.

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In particular:

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 - ▶ *Similarly for Y and Y' .*
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- ▶ ...
- ▶ *If Y is separable, we may take $Y' = Y$.*

Corollary

If Y is a separable Banach space which is a quotient of a non-separable Banach space X , then Y is also a quotient of a Banach space X' of density \aleph_1 , with $X \equiv_L X'$.

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Thank you!

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www.math.toronto.edu/cjeagle